

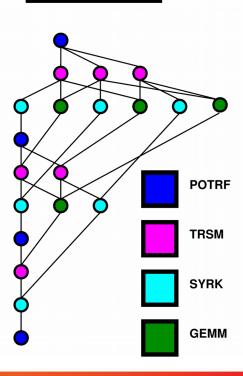
StarPU: task-based scalable Runtime system for heterogeneous multicore architectures

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Task management

Implicit task dependencies

Right-Looking Cholesky decomposition (from PLASMA)



Write your application as a task graph

Even if using a sequential-looking source code

→ Portable performance

Sequential Task Flow (STF)

- Algorithm remains the same on the long term
- Can debug the sequential version.
- Only kernels need to be rewritten
 - BLAS libraries, multi-target compilers
- Runtime will handle parallel execution



Overview of StarPU

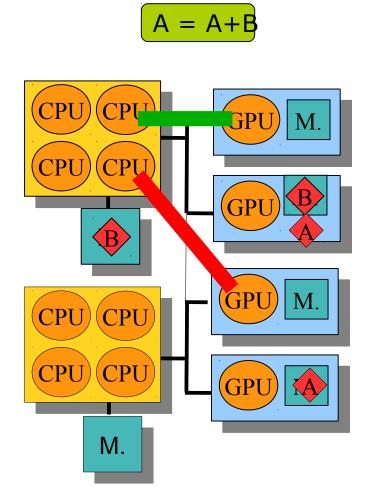
Rationale

Task scheduling

- Dynamic
- On all kinds of PU
 - General purpose
 - Accelerators/specialized

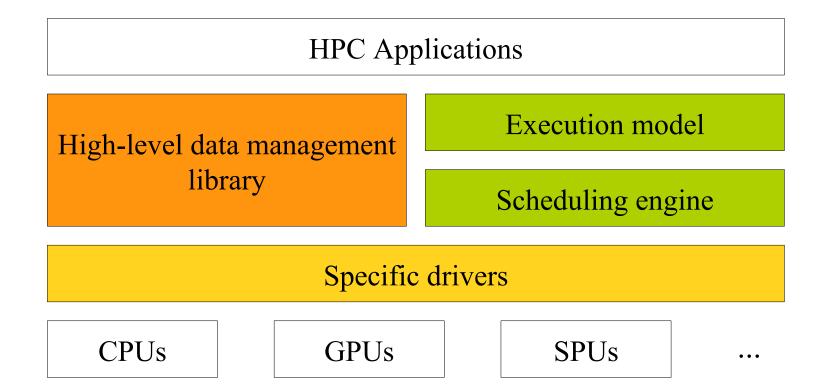
Memory transfer

- Eliminate redundant transfers
- Software VSM (Virtual Shared Memory)





The StarPU runtime system



Mastering CPUs, GPUs, SPUs ... *PUs → StarPU



The StarPU runtime system

Development context

- History
 - Started about 9 years ago
 - PhD Thesis of Cédric Augonnet
 - StarPU main core ~ 70k lines of code
 - Written in C
- Open Source
 - Released under LGPL
 - Sources freely available
 - svn repository and nightly tarballs
 - See https://starpu.gforge.inria.fr/
 - Open to external contributors
- [HPPC'08]
- [Europar'09] [CCPE'11],... >1000 citations



The StarPU runtime system

Supported platforms

- Supported architectures
 - Multicore CPUs (x86, PPC, ...)
 - NVIDIA GPUs
 - OpenCL devices (eg. AMD cards)
 - Intel Xeon Phi (MIC), Intel SCC
 - Kalray MPPA (experimental)
 - Cell processors (experimental) [SAMOS'09]
- Supported Operating Systems
 - Linux
 - Mac OS
 - Windows



Summary

```
starpu_codelet_t cl = { .cpu_func = my_f, ... };
float array[NX];
starpu_data_handle vector_handle;
starpu vector data register(&vector_handle, 0,
       array, NX, sizeof(vector[0]));
starpu task insert(&cl, vector_handle, 0);
starpu_task_wait_for_all();
starpu data unregister(vector_handle);
```



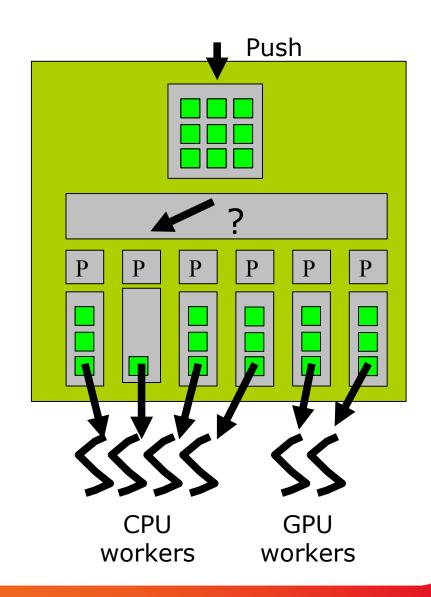
Task scheduling

Component-based schedulers

- Containers
 - Priorities
- Switches
- Side-effects (prefetch, ...)

Push/Pull mechanism

S. Archipoff, M. Sergent



More features

- Cluster support
 - MPI communication
 - Decentralized model
 - Application-provided data mapping
 - Automatic optimized transfers
- Memory consumption control
- Out of core support
 - Disk as optimized « swap »
 - or as backstore for matrix tiles
- Execution simulation support
- OpenMP and OpenCL interfaces
- ...



Applications on top of StarPU

Using CPUs, GPUs, distributed, out of core, ...

- Dense linear algebra
 - Cholesky, QR, LU, ...: Chameleon (based on Plasma/Magma)
- Sparse linear algebra
 - QR_MUMPS
 - PaStiX
- Compressed linear algebra
 - BLR, h-matrices
- Fast Multipole Method
 - ScalFMM
- Conjugate Gradient
- Other programming models: Data flow, skeletons
 - SignalPU, SkePU
- •



HiHAT wishes, in a few phrases

- Have better interface to hardware layers
 - Extremely low overhead
- Reusable components which we prefer not to maintain alone
 - Performance models, allocators, tracing, debugging, ...
- (ideally) Standard flexible task-based interface
 - Plus OpenMP etc. interfaces
 - Or at least set of helpers for outlining, marshaling, etc.
- Everything usable independently and interoperable



- Portable and flexible async APIs for driving accelerators
- Shared event management
 - Used throughout HiHAT
 - Dependencies between all kinds of requests
 - User-definable events
 - Synchronization with non-HiHAT pieces
 - Flexible event waiting API: interruptible WaitAny
 - Register a set of event
 - WaitAny(set)
 - Efficient loop over completed events
 - Can add user-defined event to the set to interrupt WaitAny easily



- Memory allocation
 - Uniform low-level API over devices
 - Efficient sub-allocator (cudaMalloc is not efficient)
 - Same-size pools
 - Allocation reuses
 - In RAM case, hierarchical balancing between cores/caches/NUMA
- Disk support : store/key/value, basically
 - store = plug(path)
 - key = allocate(store, size)
 - write(store, key, buffer, offset, size)
 - read(store, key, buffer, offset, size)
 - free(store, key)
 - unplug(store)



- Data transfer priorities
 - Take precedence over already-queued transfers
- Inter-node communication layer instead of MPI
 - Transfer priorities, again
 - Completely asynchronous and flexible event waiting API
 - Could be a PGAS: no need for messages, just memory coherency



- Performance models
 - We currently have history-based, linear reg. and multi-linear reg.
- Standard trace formats and debugger
 - We currently use paje, vite, temanejo

Could be moved to shared components

