

ARGOBOTS

A Lightweight Low-level Threading/Tasking Framework

https://collab.cels.anl.gov/display/ARGOBOTS/

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Argo Concurrency Team

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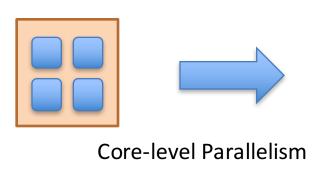
Past Team Members:

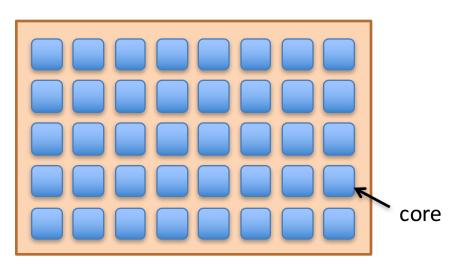
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Massive On-node Parallelism

- The number of cores is increasing
- Massive on-node parallelism is inevitable
- Existing solutions do not effectively deal with such parallelism with respect to on-node threading/tasking systems or with respect to off-node communication in the presence of such tasks/threads
- How to exploit?





Goal: Provide runtime systems utilizing massive on-node parallelism



Argo Concurrency: Approaches

High-level tasking models (CilkBots, TASCEL, PaRSEC)

- Explore high-level tasking frameworks that will take advantage of the low-level threading and communication frameworks
- CilkBots and TASCEL: PNNL & UIUC
- PaRSEC: UTK

Argobots interoperability with a PUT/GET model

- Design new communication libraries that work with such low-level threading models
- Argonne

High-level dynamic execution environments (Charm++,..)

- Investigate high-level programming models, e.g., Charm++, that are specialized in dynamic execution environments and can exploit the low-level threading and communication frameworks
- UIUC

Argobots: A lightweight low-level threading/tasking framework

- Develop a new low-level threading/tasking model that will allow us to expose the hardware characteristics of exascale computing systems more effectively
- Argonne & UIUC



Argobots

A low-level lightweight threading and tasking framework

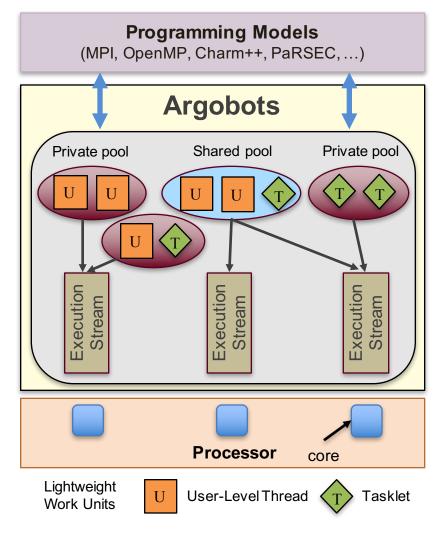
(http://collab.cels.anl.gov/display/argobots/)

Overview

- Separation of mechanisms and policies
- Massive parallelism
 - Exec. Streams guarantee progress
 - Work Units execute to completion
 - User-level threads (ULTs) vs. Tasklet
- Clearly defined memory semantics
 - Consistency domains
 - Provide Eventual Consistency
 - Software can manage consistency

Argobots Innovations

- Enabling technology, but not a policy maker
 - High-level languages/libraries such as OpenMP, Charm++ have more information about the user application (data locality, dependencies)
- Explicit model:
 - Enables dynamism, but always managed by high-level systems



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Argobots Execution Model

Execution Streams (ES)

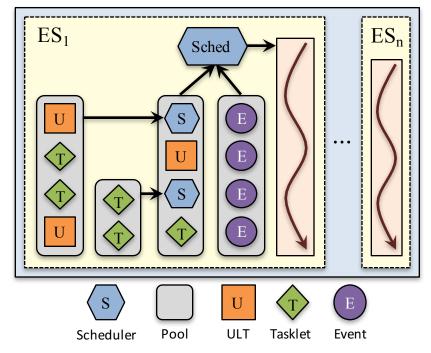
- Sequential instruction stream
 - Can consist of one or more work units
- Mapped efficiently to a hardware resource
- Implicitly managed progress semantics
 - One blocked ES cannot block other ESs

User-level Threads (ULTs)

- Independent execution units in user space
- Associated with an ES when running
- Yieldable and migratable
- Can make blocking calls

Tasklets

- Atomic units of work
- Asynchronous completion via notifications
- Not yieldable, migratable before execution
- Cannot make blocking calls



Argobots Execution Model

Scheduler

 Stackable scheduler with pluggable strategies

Synchronization primitives

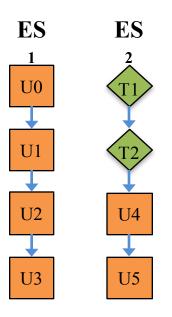
Mutex, condition variable, barrier, future

Events

Communication triggers

Explicit Mapping ULT/Tasklet to ES

- The user needs to map work units to ESs
- No smart scheduling, no work-stealing unless the user wants to use



- Benefits
 - Allow locality optimization
 - Execute work units on the same ES
 - No expensive lock is needed between ULTs on the same ES
 - They do not run concurrently
 - A flag is enough

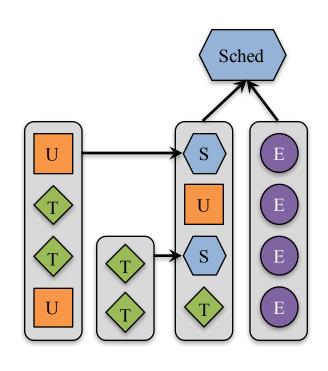
Optional Advanced features: shared and global workpools



Stackable Scheduler with Pluggable Strategies

- Associated with an ES
- Can handle ULTs and tasklets
- Can handle schedulers
 - Allows to stack schedulers hierarchically
- Can handle asynchronous events
- Users can write schedulers
 - Provides mechanisms, not policies
 - Replace the default scheduler
 - E.g., FIFO, LIFO, Priority Queue, etc.
- ULT can explicitly yield to another ULT
 - Avoid scheduler overhead



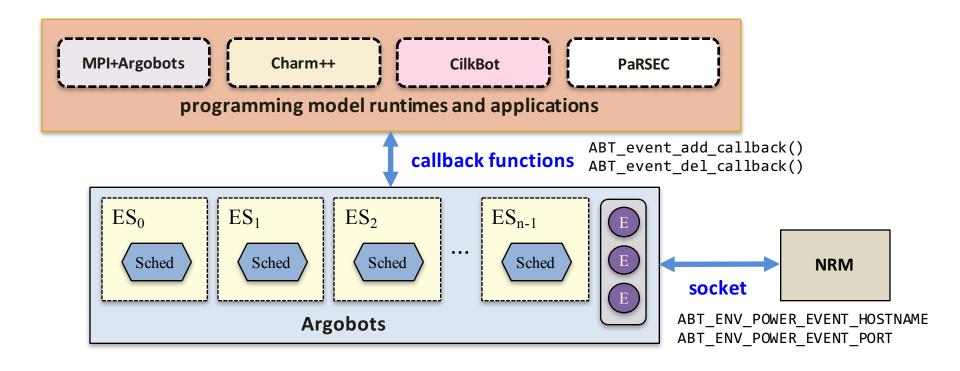


Some Capabilities of Argobots

- Shrinking the set of ESs used
 - E.g. for power/energy optimization
- Expanding the set of ESs used
- Migration of tasks and ULTs
 - Load balancing
 - Locality
 - NUMA issues



Interfaces for Shrink/Expand Events

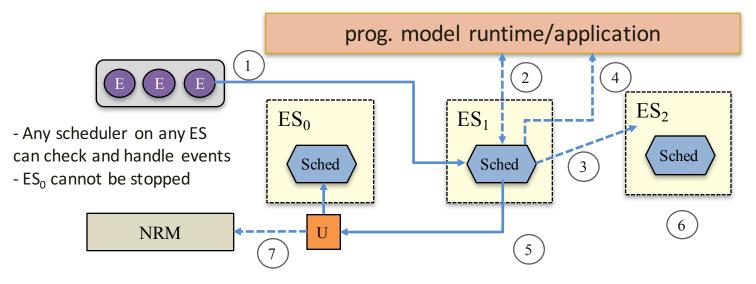


- [Argobots] Connect to NRM using a socket on ABT_init()
- 2. [Runtimes/applications] Register callback functions for shrink/expand events
- 3. [Runtimes/applications] Deregister callback functions when they terminate
- [Argobots] Disconnect from NRM on ABT_finalize()



Shrink/Expand Event Handling

Shrinking

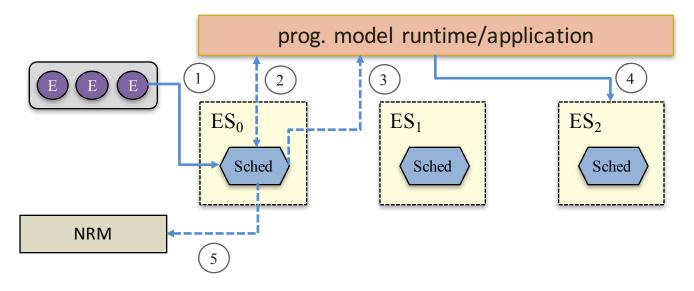


- 1. ES₁ picks an event, which requests ES₂ to be stopped
- 2. Ask the runtime using callbacks whether ES₂ can be stopped
- 3. If OK, mark ES₂ to need to stop so when the scheduler on ES₂ checks events, it can be stopped
- 4. Notify the runtime that ES₂ will be stopped
- 5. Create a ULT on ES_0
- 6. When the scheduler on ES₂ stops, ES₂ is terminated
- 7. After ES₂ is terminated, the ULT frees ES₂ and sends a response to NRM



Shrink/Expand Event Handling

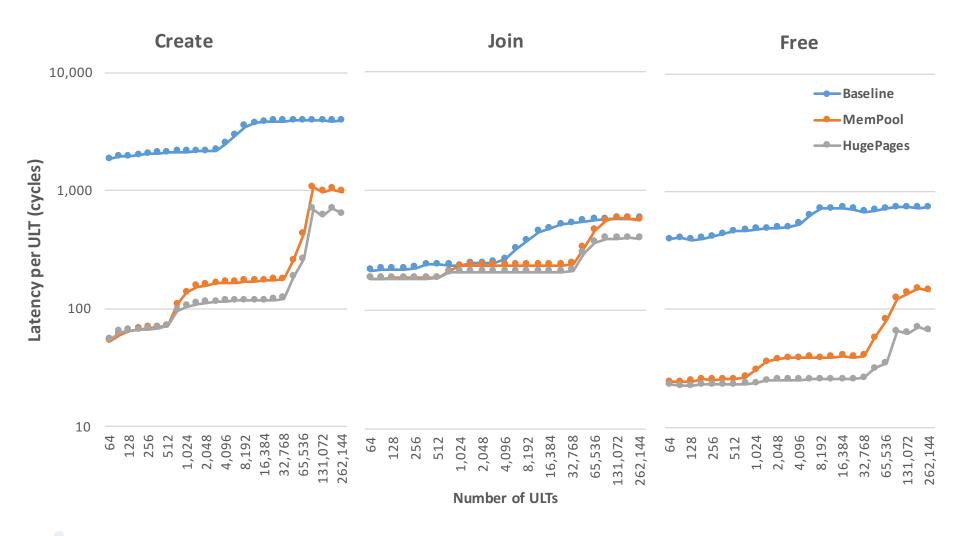
Expanding



- 1. ES₀ picks an event, which requests to create an ES₂
- 2. Ask the runtime using callbacks whether it can create ES₂
- 3. If OK, invoke a callback function so the runtime creates ES₂
- 4. Create ES₂
- 5. Send a response to NRM

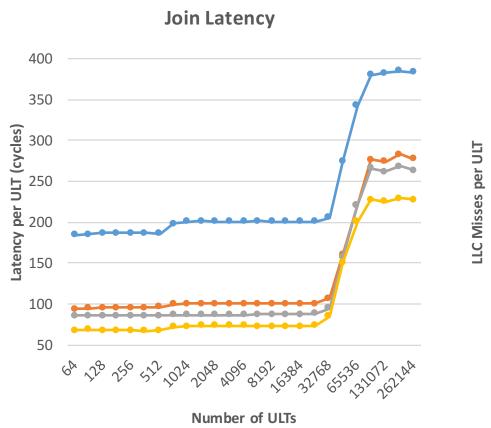


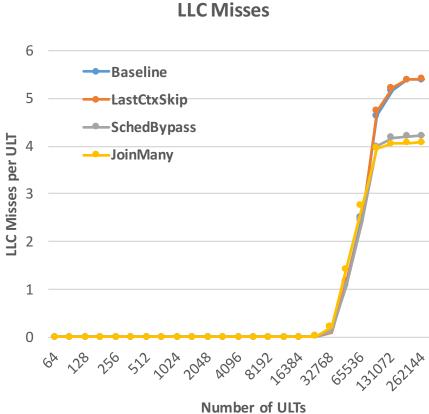
Optimization: Memory Pool & Huge Pages



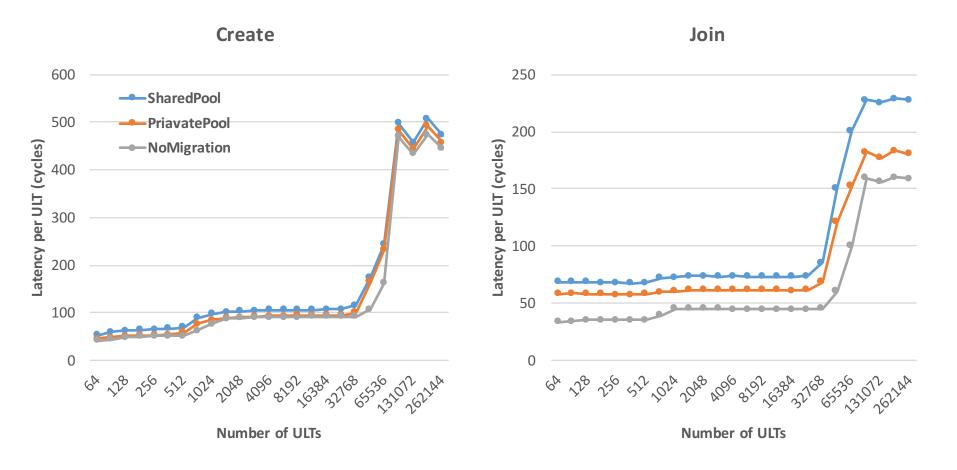


Optimization: Join Operation



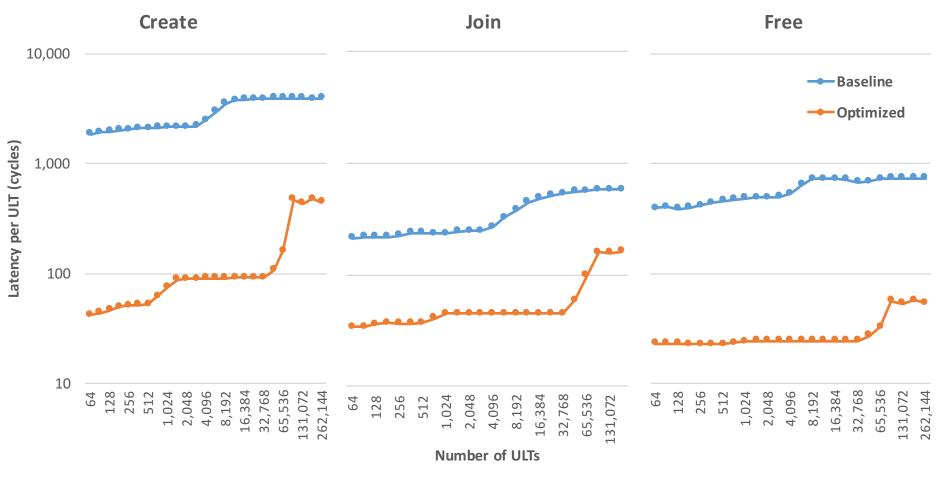


Optimization: Private Pool and Disabling Features





Optimization: Putting It All Together



- 64 ULTs : 2,443 cycles (c: 1,837, j: 212, f: 394) -> 99 cycles (c: 42, j: 34, f: 23)
- 262,144 ULTs: 5,212 cycles (c: 3,921, j: 558, f: 733) -> 659 cycles (c: 446, j: 159, f: 54)

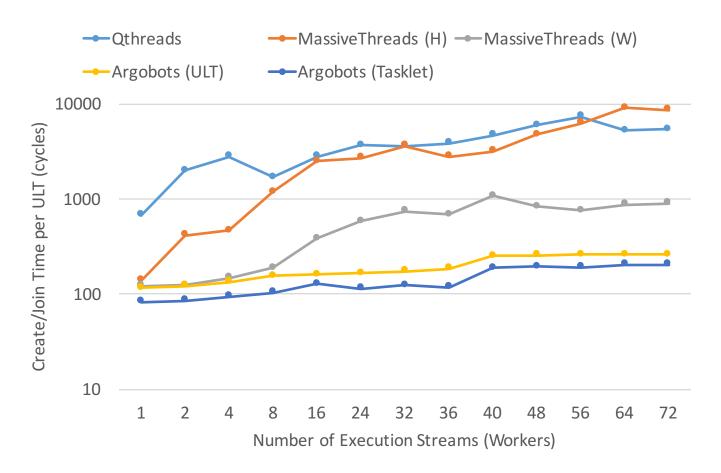
24.7x speedup

7.9x speedup



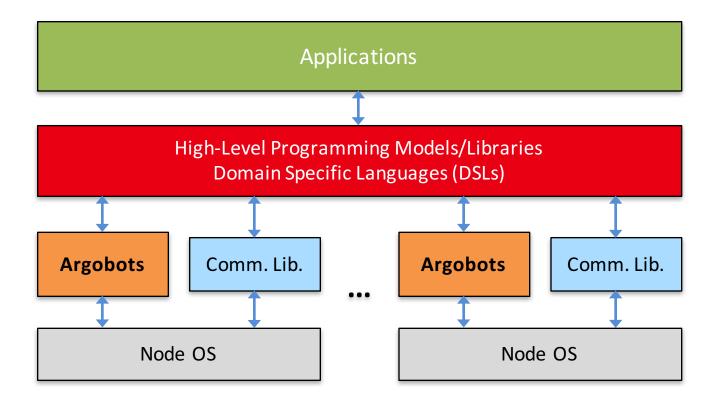
Performance: Create/Join Time

- Ideal scalability
 - If the ULT runtime is perfectly scalable, the time should be the same regardless of the number of ESs



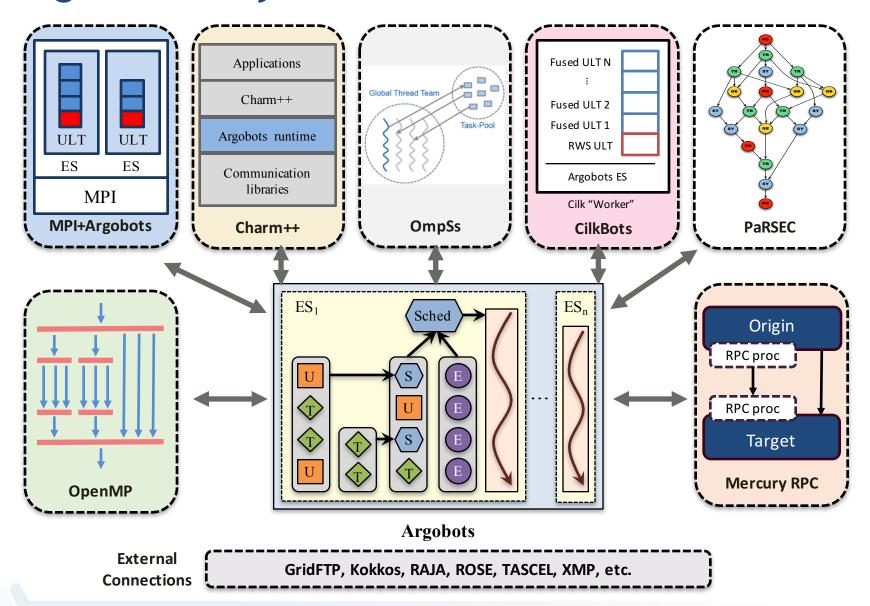


Argobots' Position



Argobots is a low-level threading/tasking runtime!

Argobots Ecosystem



Argobots-Aware MPI Runtime

Problem

- Traditional MPI implementations are only aware of kernel threads
- Thread-synchronization costly to ensure thread-safety and progress requirement from MPI
- Wasted resources if a kernel thread blocks for MPI communication

Solution

- An MPI implementation aware of Argobots threads
- Lightweight context switching to overlap costly blocking operations (communication, locks, etc.)
- Reduced thread-synchronization opportunities (guaranteed consistency within an ES without locks or memory barriers)

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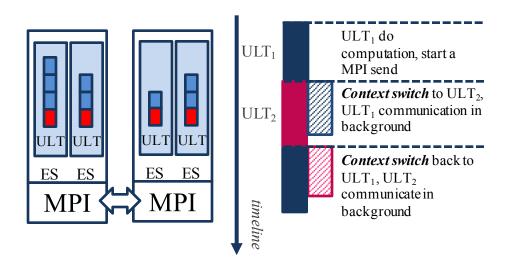
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Recent results

- Developed an MPICH+Argobots prototype
- Demonstrated the ability to overlap blocking communication with HPCG, SpMV, etc.
- Deployed successfully a fully threaded Graph500 benchmark implementation

Impact/Potential

 The new MPI+Argobots model has the potential to overcome the long lasting multithreaded MPI communication challenge

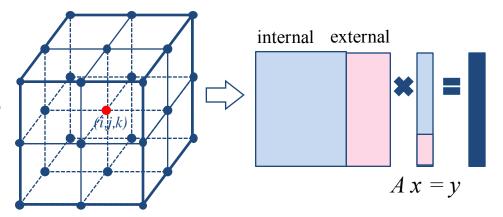


MPI+Argobots Execution Model



Application: HPCG

- High Performance Conjugate Gradient (HPCG)
 - Solves Ax=b, large and sparse matrix
- Hiding Global Collective Communication
 - at end of each iteration, do
 DDOT to calculate tolerance (vector multiplication + MPI_Allreduce)
 - overlap communication and computation (overlap DDOT with MG in the next iteration)
 - fork a ULT to do ult_ddot and join in the next iteration

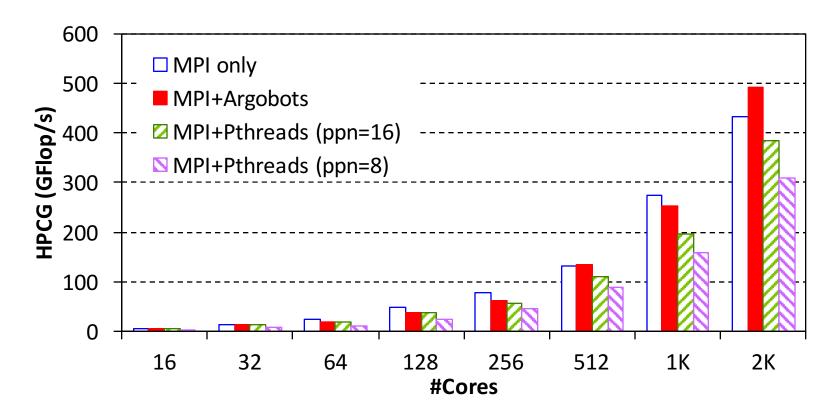


MG: preconditioner of CG

DDOT: dot product follow by MPI_Allreduce



Preliminary Results: HPCG



- On 2,048 cores, HPCG using MPI+Argobots shows performance improvement of 13.4% over MPI-only version, or 27.4% over MPI+Pthreads version.
 - As core number increases, the benefit of communication hiding begins to reveal. DDOT% increases from 0.62% on 16 cores to 36.8% on 2,048 cores.



OpenMP over Argobots

Problem

- OS thread-based OpenMP implementations may not be efficient for massive dynamic parallel applications
- A blocking communication call makes an entire execution stream (i.e., OS thread) blocked
- Inefficient support of nested parallelism and tasks

Solution

- Develop an OpenMP compiler that generates Argobots ULTs and tasklets depending on the existence of blocking call (e.g. MPI call) in the code
- Develop an OpenMP runtime that manages created ULTs and tasklets over a fixed set of computational resources

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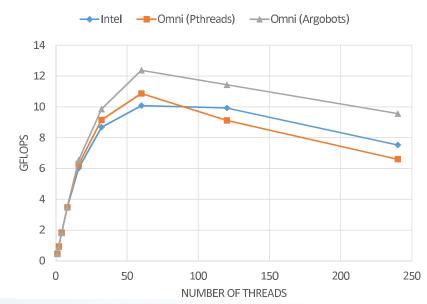
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Recent results

- Developed an OpenMP compiler/runtime prototype over Argobots based on Clang/LLVM (ANL): BOLT compiler and runtime framework
- Optimized Omni compiler framework (collaboration with RIKEN)
- Demonstrated better support of nested parallelism

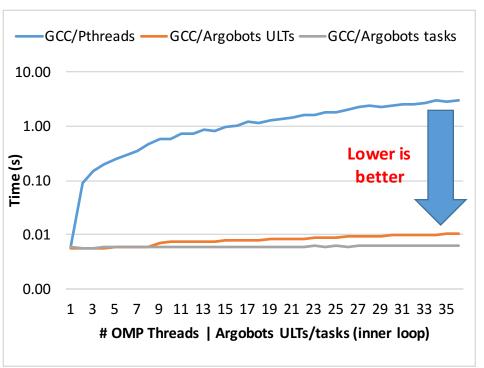
Impact/Potential

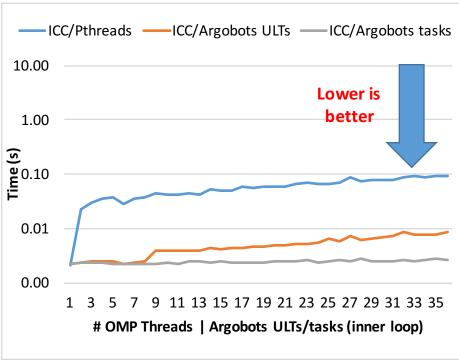
 The new OpenMP compiler/runtime developed would improve the performance of applications by leveraging lightweight threads/tasks



Nested Parallel Loop: Performance

Execution time for 36 threads in the outer loop





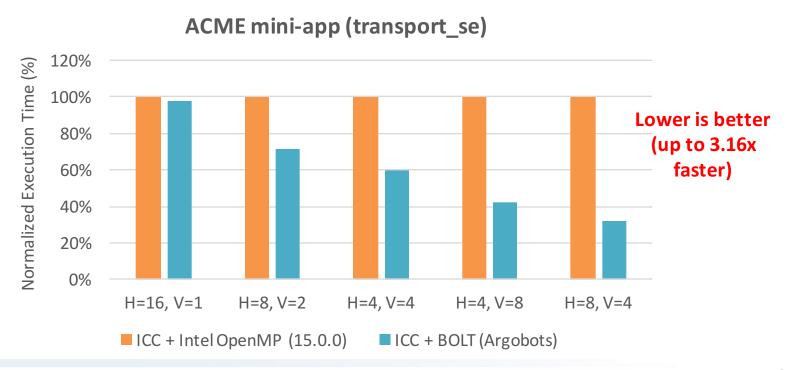
GCC OpenMP implementation does not reuse idle threads in nested parallel regions, all the teams of threads need to be created in each iteration

Some overhead is added by creating ULTs instead of tasks



Application Study: ACME mini-app

- ACME (Accelerated Climate Modeling for Energy)
 - Implementing additional levels of parallelism through OpenMP nested parallel loops for upcoming many-core machines
- Preliminary results of testing the transport_se mini-app version of HOMME (ACME's CAM-SE dycore)



Charm++ with Argobots

Problem

- Charm++ is a very rich parallel programming model and several important scientific applications • (like NAMD, OpenAtom, ChaNGA etc.) have been written using it.
- In this project, our focus is to develop an Argobots port for Charm++.
- We also address the problem of dynamically shrinking/expanding the number of cores for better power management.

Argobots Port for Charm++

- Converse is the active messaging layer in Charm++.
- We prepare the Argobots port for Charm++ by instantiating Charm++ pthreads as Argobots execution streams.
- Instead of enqueing Charm++ messages in Converse queues, we put them into appropriate Argobots pools as Argobots tasks.
- Finally, we implement Charm++ threads on top of Argobots threads with condition variables to implement suspend/resume.

Charm++ model

Intelligent runtime

Converse runtime

(threading, messaging, scheduler)

Argobots
(ULTs, Tasks, Scheduling, etc.)

Communication libraries (MPI, uGNI, PAMI, Verbs)

Shrink/Expand

- Charm++ maintains a set of pools for each scheduler, mapped to an xstream
- Ranks in Charm++ are virtualized by saving the mapping of pool to a xstream
- When an xstream is removed, the associated pools are put into a global list
- To maintain correctness in Charm++, the rank of any tasks/threads in the global list are derived from the pool (ranks are virtualized)
- Shrink: Other xstreams execute work units from orphaned pools with some added synchronization
- Expand: A new xstream is created and takes over a set of orphaned pools Contact:
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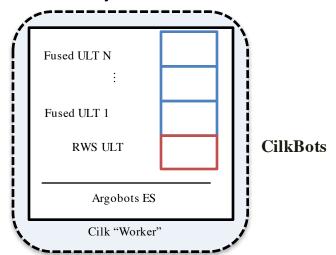
CilkBots: Lightweight Threading-Based Cilk

Problem

- Massive on-node parallelism and memory domains
- Existing spawn-sync tasking frameworks are often not inter-operable and do not account for data locality

Solution

- Cilk runtime that exploits Argobots low-level threading/tasking model
- Fused execution of user-level threads in a single
 Cilk worker. Constrained interleaving to maximize cache locality



Recent results

- CilkBots prototype using Argobots
- Demonstrated cache locality optimization for CilkBots programs
- Demonstrated ability to dynamically change the number of ESs to respond to external events (e.g., power management)
 - Dynamically rebalances and executes to completion around changes in available resources

Impact/Potential

 CilkBots allows design of spawn-sync programs that are responsive to external events, inter-operate with other models, and enable dynamic locality optimization

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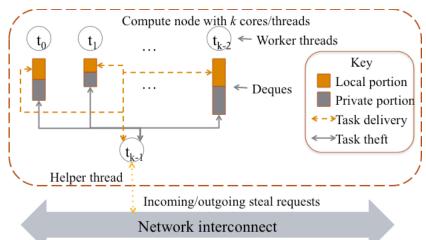
TASCEL with Argobots

Problem

- To maximize performance, distributed-memory tasking runtimes require significant programmer effort to expose all blocking operations
- These runtimes also take over program control and do not inter-operate with other runtimes in a program's execution

Solution

- Distributed-memory work stealing using lightweight threads (ULTs)
- Context switching between ULTs for performance and inter-operability



Recent results

- TASCEL prototype using Argobots
- Demonstrated ability to dynamically change the number of ESs to respond to external events (e.g., power management)

Impact/Potential

 TASCEL over Argobots enables the design of more performance portable and composable task-parallel program phases

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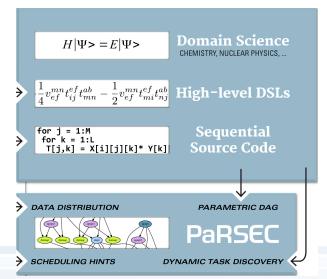
PaRSEC: A distributed runtime for domain specific languages

Problem

- Performance portability, a critical component of HPC, is difficult to be guaranteed by programming paradigms on Complex heterogeneous distributed architectures
- Algorithmic advances and architectural progress are often disjoint

Solution

- separation of concerns: compiler optimizes tasks, developer describes dependencies between tasks, runtime orchestrates the dynamic execution
- Interface with the application developers through specialized domain specific languages
- Separate algorithms from data distribution



Recent results

- Full integration of PaRSEC with Argobots
- Full support for distributed executions
- Initial support for marshaling resources (e.g., accelerators) from NodeoOS
- Add support for task stealing and migration to adapt to dynamic condition (power management, failure prediction)
- Application support: Support for the ScaLAPACK library, and initial integration with NWCHEM

Impact/Potential

- Remove all control flow from application, resulting on completely data-driven executions
- Ease the user adoption by providing rich and flexible domain specific languages

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Argo Concurrency: Recent Results & Impact

- Recent results
 - Initial development of the "Argobots" threading/tasking library
 - Initial design of extending inter-node communication in the presence of user-level threads and tasks
 - Prototype development of OpenMP, MPI, Charm++,
 CilkBots, TASCEL, and PaRSEC over Argobots
- Impact
 - The new models developed would impact both legacy and new programming models that take advantage of massive on-node parallelism

